

Java Access to CA-IDMS Data at BT Session Abstract

 BT, one of the world's leading providers of communications solutions and services, has recently introduced direct access to its major CA-IDMS CSS (Customer Service System) databases using the CA-IDMS Server JDBC interface.



- This session explains why, how, and a few of the lessons learnt so far and the progress made.
- Steve Rundle has worked for BT for 40 years, most of these years as an IDMS DBA. He has been the key DBA for the IDMS access via java project (JCSS) for the last year.



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Java Access to CA-IDMS Data at BT Topics

- · IDMS in BT
 - Mainframe Modernisation Program
 - JCSS IDMS Gateway Overview
 - CA-IDMS JDBC Configuration
 - Suspend Strategy / Active Tasks
- · What is the SQL statement actually doing?
- Dynamic SQL CACHE
- JCSS Current Status 2018
- Performance and measurement
- SQL Statement Validation
- · SQL Table procedures
- · Virtual Keys Schema
- SQL Update
- · Two Phase Commit
- · Schema Definitions
- Q&A





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Java Access to CA-IDMS Data at BT IDMS in BT

What is BT

- British Telecommunications
 - UKs largest Telecomm provider
 - 4 tiers of retail product
 - Telephone/Internet/Mobile/TV
 - Business Network provision and management
 - System management
 - Global Services
 - 60 Million Customers
 - >50,000 employees worldwide





Java Access to CA-IDMS Data at BT IDMS in BT

- 4 Major CA-IDMS Applications at BT (CAMSS,CSS,TS/OPS & EXPRESS)
- CSS Customer Service System is the largest
 - Holds customer info, and Local provision of service details
 - · 30 years old
 - Split Into 29 Regions of the UK each running on its own LPAR
 - Each region has one Update, one Read Only (using DB-EZ/Synchro) and a maintenance CV
 - CICS Front End / CA IDMS Database
 - Bespoke transaction switching solution appears as a single DB
 - Application is predominately COBOL
 - Bespoke Common Middleware solution provides application services





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Java Access to CA-IDMS Data at BT Mainframe Modernisation Program

- BT Recently undertook a review of the Mainframe Platform
 - Driven by Cost Reduction, Sustainability, System Rationalisation
 - Expertise and experience being lost
 - Re-platforming of all mainframe applications was explored
 - Outcome was to retain the platform, seek to reduce costs, address sustainability issues and modernise where possible to become a 'better fit' with the rest of the BT estate
 - Several proof of concepts being undertaken including enabling java access to CA-IDMS to provide a CSS 'data as a service' interface running under Linux on a separate platform





Java Access to CA-IDMS Data at BT JCSS IDMS Gateway Overview

- Java access to CA-IDMS to provide a CSS 'data as a service' (Linux on a separate platform)
- CA-Server JDBC to/from the 29x2 CSS IDMS databases
- · Business logic replicated in java
- · Reduce CPU processing costs
- · Business Logic
 - COBOL to Java
 - Knowledge retention





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Java Access to CA-IDMS Data at BT JCSS IDMS Gateway Overview

- Java access to CSS "JCSS"
 - Business logic replicated in Java
 - Native SQL to be used
 - · Even where data is redefined
 - IDMS R19 Virtual Foreign Key feature (esp. for update)
 - Requires some table Procedures
 - · ODO and Redefined CALC keys
 - Requires some schema changes





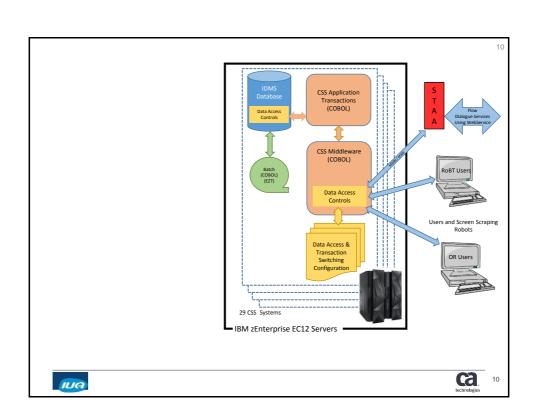
Java Access to CA-IDMS Data at BT JCSS IDMS Gateway Overview

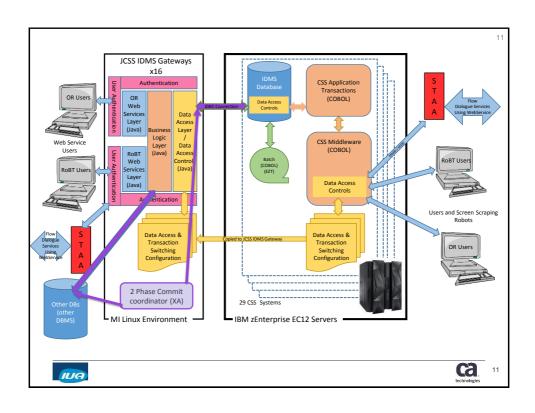
- Native SQL to be used whenever possible to avoid the overhead of maintaining table procedures
- We will be using the IDMS R19 Virtual Foreign Key feature
- Requires some schema changes but not material changes
 for example
 - FILLER redefined with alternative element name
 - Reordering redefined elements within a record
 where redefined elements are being used as set keys. We need to reorder the definitions / redefinitions to make the key values the base elements
 - · Defining OM Indexes as MA.

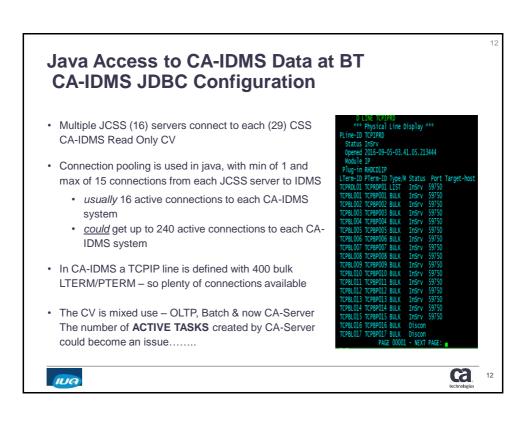




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Java Access to CA-IDMS Data at BT Suspend Strategy / Active Tasks

- Suspend Strategy for Pseudo-Conversational Processing has an impact.
- Controls when CA-IDMS Server will issue a suspend which ends the CA-IDMS task and frees resources on the CV
- This behaviour can be customised by selecting a "suspend strategy"
- Pre-defined set of strategies available appropriate for particular types of application.
- Changed from INTERACTIVE to SERVICE
 - Concern about MAX TASKS
 - · Higher Storage





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Java Access to CA-IDMS Data at BT Suspend Strategy / Active Tasks

 Interactive—Intended for use by applications with a user interface, in which database activity and user input are intermixed.

Task remains active on connect. Suspends when the transaction is committed. This is the default strategy for ODBC and non-pooled JDBC connections.

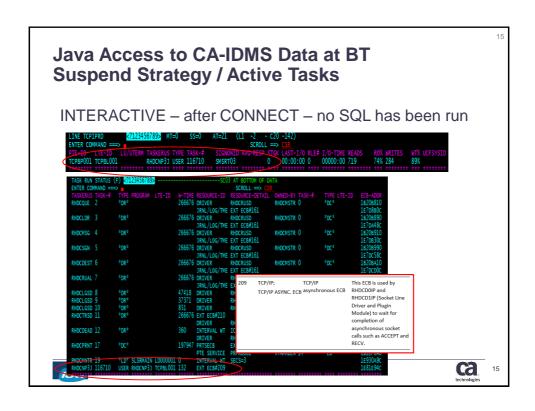
Service—Intended for use by JDBC applications that run in an application server that
pools connections and allocates them temporarily to units of work that that access the
database one or more times without waiting for user input.

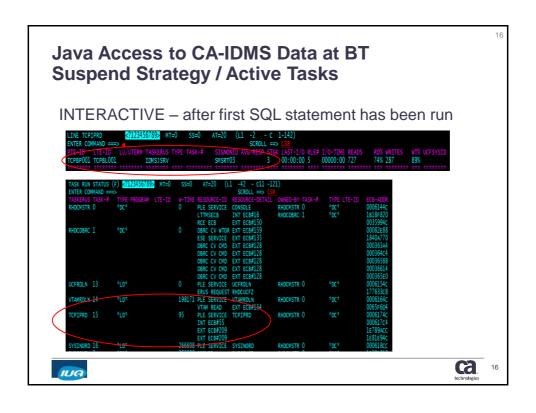
Task remains active on commit – until it determines the connection is idle. This is the default for pooled JDBC connections.

- Batch—Intended for use by applications access the database many times and terminate without waiting for user input. The driver does not use pseudo-conversational processing at all.
- Custom









Java Access to CA-IDMS Data at BT

- What is the SQL statement doing in IDMS?
- Service affecting resource hungry SQL
- How do we know?

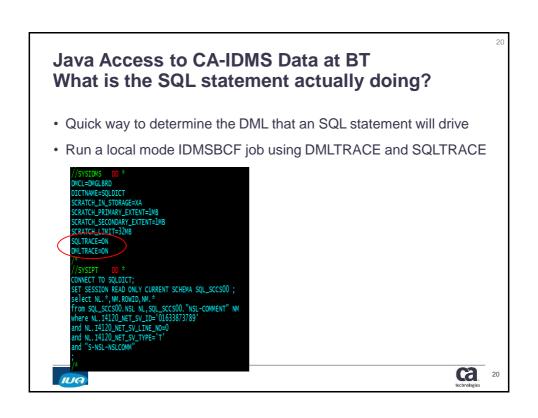
Explain
DML Trace
Experience

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Java Access to CA-IDMS Data at RT CHAR(18) Set or index name What is the SQL sta CHAR(1) Leaf scan indicator, when ACMODE is I. This indicates whether data is retrie by sequential access to index leaf pages. 'N'-No · EXPLAIN can be run for tl Composite sort type. A nonblank value in this field indicates an actual sort is Inter EXPLAIN STATEMENT 'SEL WHERE NL.14120_NET_SV_1 required (data cannot be accessed in sort order). type o 8-DE Mode of access to the database underlying the table, when STYP 'G'—Group 'A'—Area 'M'—Merge join 'C'-CALC 'O'—Order by 'I'—Index Inner sort type. This is an actual sort performed for the inner loop of a merge CHAR(1) 'M'—Set member 'O'-Set owne 'M'—Merge join PDICT TABLE 'P'-Table procedure SUBQC CHAR(1) Subquery correlation. 'S'—Sequential 'T'—(Temporary table) 'N'-Not correlated 'Y'—Correlated 0001-01-01-00.00.00.00 0001-01-01-00.00.00 2016-09-20-10.51.19.117 Ca IUA

```
Java Access to CA-IDMS Data at BT
Option ===>
Press PF6 to view SQL Syntax File
SQL: SELECT NSL.ROWID, NSL.*, ROU
                                                                           SCHEMA:
SCHEMA:
TABLE:
ACCESS MODE:
SET OR INDEX:
      DBNAME:
STEP:
COMMAND:
STEP TYPE:
PARENT BLOCK:
PARENT STEP:
STEP:
COMMAND:
STEP TYPE:
                                                      COMPOSITE SORT:
INNER SORT:
                                                                                                                            IX LEAF SCAN: SUB QUERY CORR:
                                                                           SORT:
SCHEMA:
TABLE:
ACCESS MODE:
SET OR INDEX:
 QUERY BLOCK:
PARENT BLOCK:
PARENT STEP:
STEP:
COMMAND:
STEP TYPE:
                                                      COMPOSITE SORT: INNER SORT:
                                                                                                                            IX LEAF SCAN: SUB QUERY CORR:
                                                                           SCHEMA:
TABLE:
ACCESS MODE:
SET OR INDEX:
 QUERY BLOCK: 0
PARENT BLOCK:
PARENT STEP:
                                                      COMPOSITE SORT:
INNER SORT:
                                                                                                                            IX LEAF SCAN: SUB QUERY CORR:
                                                                                                                                               6
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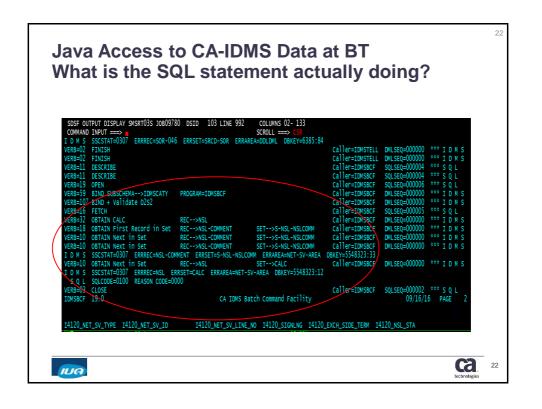


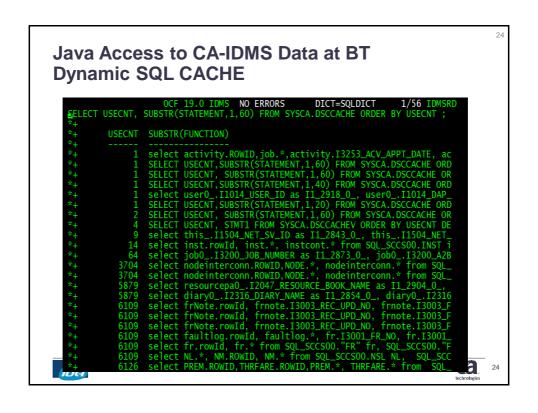
```
Java Access to CA-IDMS Data at BT
What is the SQL statement actually doing?
                         DISPLAY SMSRT03S JOB09780 DSID 103 LINE 24
       SDSF OUTPUT DISPLAY SUSRIUS JOBUS/NO DSID 103 LINE 24 COLUM
COMMAND INPUT ==> SCROLL
STATUS = 0 SQLSTATE = 00000
SEF SESSION READ ONLY CURRENT SCHEMA SQL_SCCS00;
SEF SESSION READ ONLY CURRENT SCHEMA SQL_SCCS00
STATUS = 0 SQLSTATE = 00000
SELECT NL.* NN.ROWID, NN.*
SELECT NL.* NN.ROWID, NN.*
                                                                                                                                     Caller=IDMSBCF SQLSEQ=000013 *** S Q L
            m SQL_SCCSOO.NSL NL, SQL_SCCSOO."NSL-COMMENT" NM
re NL.14120_NET_SV_ID='01633873789'
NL.14120_NET_SV_LINE_NO=0
       and "S-NSL-NSLCOMM

VERB=20 PREPARE--> select NL.*,NM.ROWID.NM.*

from SQL_SCCS00.NSL NL.SQL_SCCS00."NSL-C

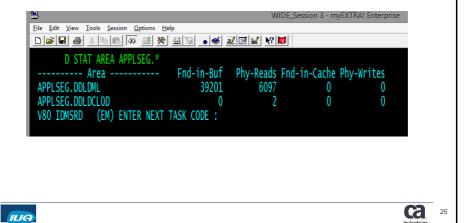
where NL.I4120_NET_S
                                                                                                                                      Caller=IDMSBCF SQLSEQ=000007 *** S Q L
                                  OMMENT" NM
V_ID='01633873789'
                                  OMMEN: 18-3
V_ID="0163873789"
and NL.14120_NET_SV_LINE_NO=0
and NL.14120_NET_SV_TYPE="T"
and "S-NSL-NSLCOMM"
       VER8=59 BIND SUBSCHEMA-->IDMSCATZ PROGRAM=IDMSDDAM
VER8=37 READY Area Retrieval AREA->DDLCAT
                                                                                                                                      Caller=IDMSDDAM DMLSEQ=000000 *** I D M S Caller=IDMSTELL DMLSEQ=000000 *** I D M S Caller=IDMSTELL DMLSEQ=000000 *** I D M S
                  READY Area Retrieval
BIND Record
                                                             AREA->DDLCAT
REC-->AREA
                                                                                                 ADDR=9B6D0E10
ADDR=9B6D0ECC
                                                                                                                                                                                                   ca
                                                                                                                                                                                                                   21
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```





Java Access to CA-IDMS Data at BT Dynamic SQL CACHE

 Keep an eye on I/O for the dictionary – if this rises it could indicate that the SQL cache is not effective



Java Access to CA-IDMS Data at BT JCSS - Current Status 2018

- 22 Transactions now developed and executing in production
- 37 Million Java Tasks executed per day (Not the same as Transactions - Approx 3 Java tasks per Xact)
- IDMS Update SQL analysed
- Monitoring and Measurement challenges
- Suspend Strategy to be changed to SERVICE
- Table Procedures developed Overcome IDMS SQL limitations
- 2 –phase commit possible



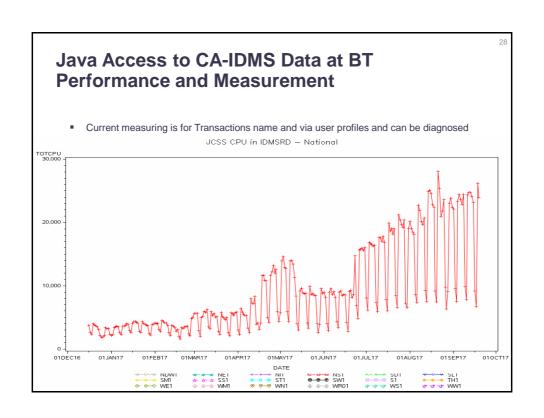


Java Access to CA-IDMS Data at BT Performance and Measurement

- · Measured Detailed CICS transactions
 - Type/ Userid/ Usage/ Business Usage
 - We record and report on CPU/Storage/IOS for each one.
- · JCSS we can only record JSRV TASKS
 - Not granular, only time and Place
- · But we are looking for further details
 - JCSS Server records Transactions/type/user etc
 - We can look at SQL Cache for types of SQL queries
 - CA are looking at an API change

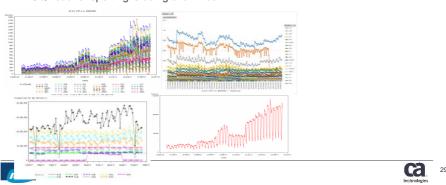
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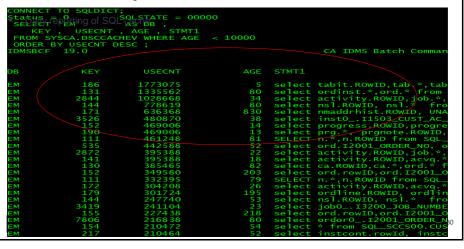
Java Access to CA-IDMS Data at BT Performance and Measurement

- Current measuring is for Transactions name and via user profiles and can be diagnosed to individual users, But java tasks are not so easily monitored
- We are "flying blind" with only headline total Xactions and CPU usage to manage.
- We are looking to Coding the JAVA with a program name. So allow IDMS to use the program name
- SQL cache reporting is being examined.



Java Access to CA-IDMS Data at BT Performance and Measurement SQL Cache

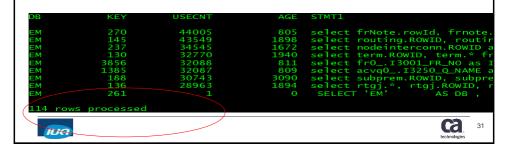
- Careful reporting on SQL Cache is possible to identify high quantity SQL queries
- SQL cache maintained against EACH of the 29 CVs.



Java Access to CA-IDMS Data at BT Performance and Measurement SQL Cache

- Keeping an eye on these can show what SQL is being used most.
- Not easily aligned with business transactions because SQL Statements are re-used across Transactions. But still a good indicator of usage.
- We also keep an eye on the capacity (number of most recent frequently used Statements, (Selected by AGE)
- We use 200 CACHE entries. If it ever goes above this usage for frequently used, recent statements then performance will suffer.

EG. SELECT from DSCCACHEV where age < 5000 and usecnt > 10000

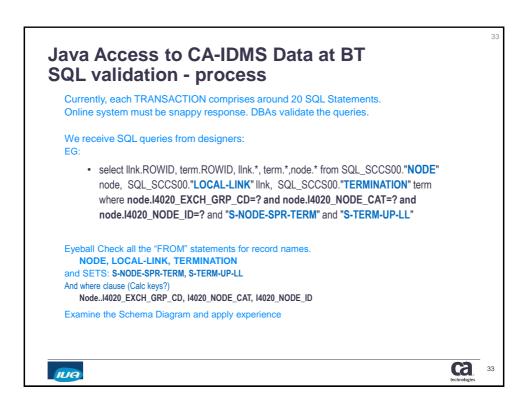


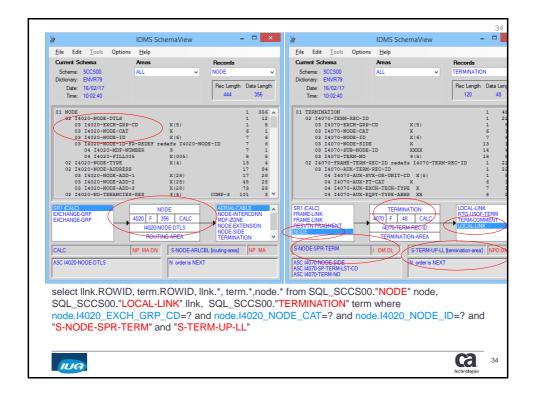
Java Access to CA-IDMS Data at BT SQL Statement validation

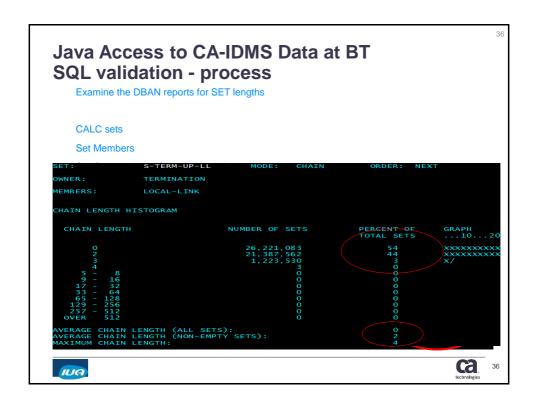
- Validate Java SQL generated Queries for efficiency.
 Educate designers to IDMS
- SQL Queries are submitted to validation and authorisation
- Three processes
 - Eyeball the SQL
 - EXPLAIN
 - DML Trace

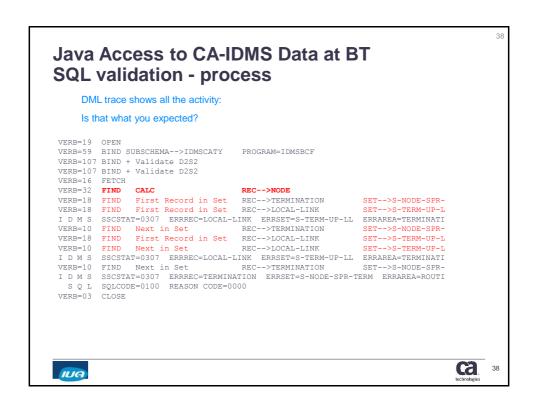


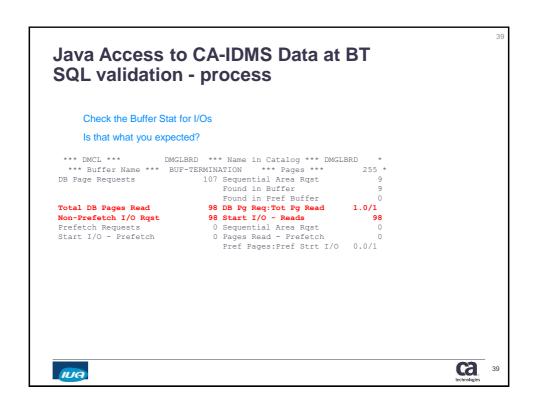


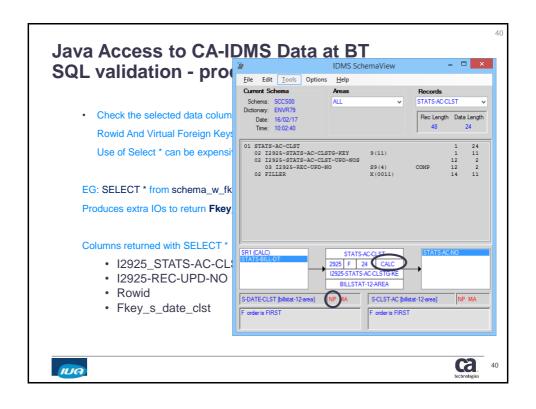


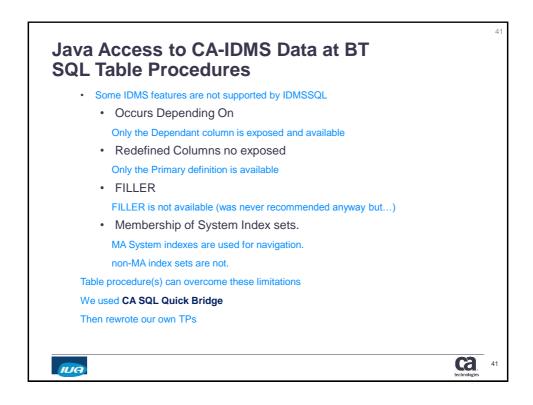


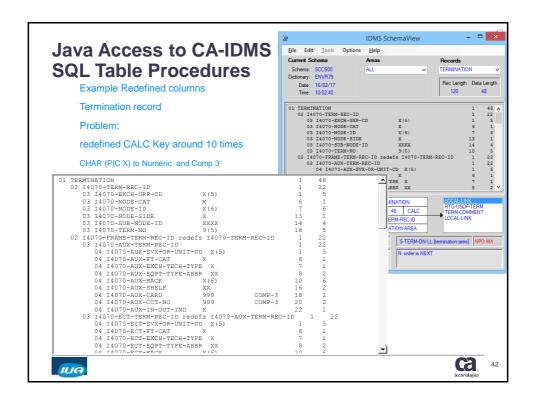


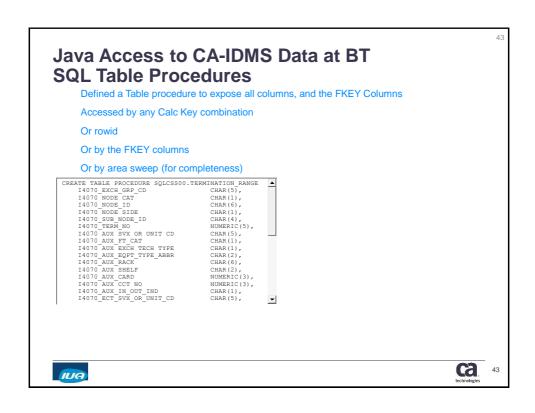


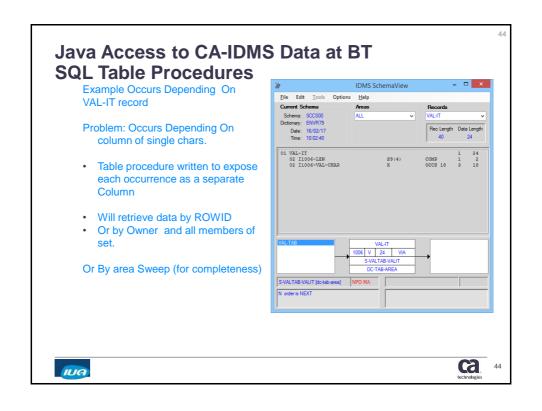


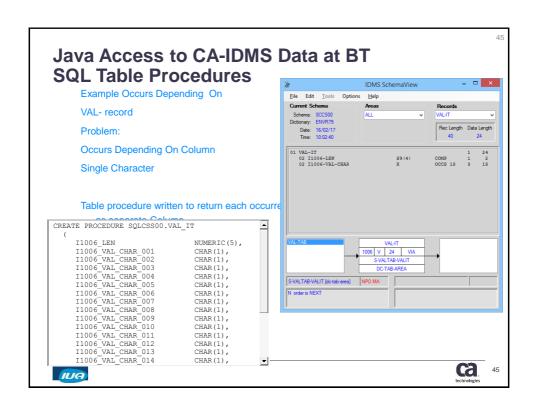


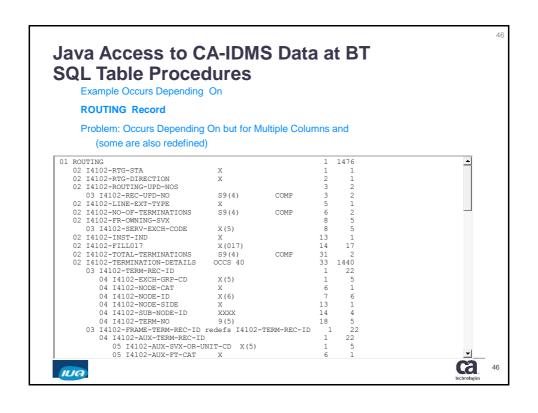


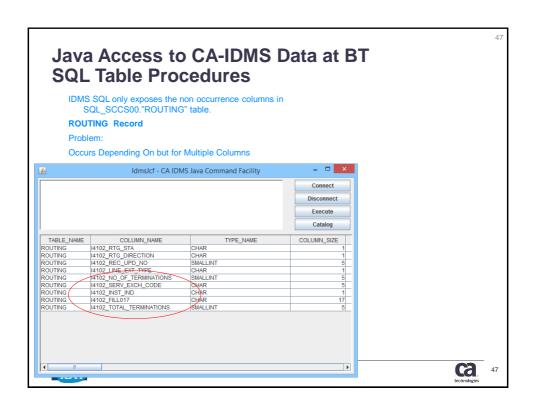


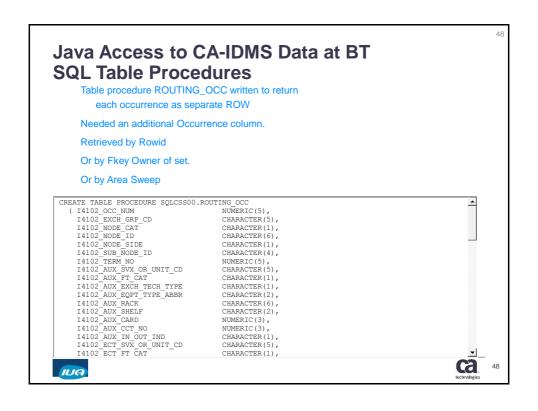












Java Access to CA-IDMS Data at BT SQL Table Procedures

Things to watch with TP:

WHERE Clause columns only passed to TP if "=" is specified

e.g. WHERE column > x is not seen by Table Procedure.

Whatever is supplied in WHERE column = must be returned exactly

eg: don't be clever with data

Code with MODE = BATCH unless DC functions required

TP program is called at least 3 times in each query.

- 1- Initiate
- 2- Called each time to return a row
- 2a return with no-more-rows
- 4 Finish

We Treat TPs like Database tables, I.E. all data requests are valid

EG: cater for daft SQL requests and Area sweeps anyway

TP call does not pass columns in SELECT, So return all columns
Care taken to return FKEY columns. Could result in Multiple I/Os
We have used "WHERE FKEY = 0" to tell TP don't return FKEY column.





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Java Access to CA-IDMS Data at BT VIRTUAL KEYS Schema

Existing SQL may not work as expected

SELECT rowid, * from tablename returns all data columns and rowid

But against VFK Schema will return two ROWID columns, plus additional FKEY columns.

If Member records do not hold Owner Pointers, satisfying the implied request for OWNER will result in IDMS reading the set to the Owner.

Many User Owned Index sets do not hold Owner pointers. So inadvertently an SQL query one could result in additional I/Os. extra I/O is expensive. Also MULTIPLE owners

- · Rowid is accepted differently in the two Schemas
- EG: select xxx where rowid = x'0FB14401' Works on NON-FK schema
- · Does not work on VFK schema





Java Access to CA-IDMS Data at BT **SQL Update**

- Update transactions under way.
- Concerns are:
 - · Locking is assumed to be "optimistic".
 - Retrievals are not locked before updates
 - SQL update process differs and Locking may introduce more Deadlocks
 - Journals can produce new record types
 - · Not all updates are possible
 - Table procedures need to be capable of update
 - 2 phase commit
- Set connections require the latest "WITH VIRTUAL KEYS" schema





Java Access to CA-IDMS Data at BT **SQL Update**

Journal records

RTSV: Return to Save Point

Generated because SQL needs to roll back to a point not marked with BGIN.

Testing has taken a lot of time:

IDMS handles this JRNL record perfectly. Rollback and RFWD and Warmstart all work perfectly

Some of our other utilities need amending to cope.





Java Access to CA-IDMS Data at BT **SQL Update** RTSV example **EG: This Transaction** UPDATE Account Set ac_balance = 123.56 where ac_num = 1234; (this works ok) **UPDATE Invoice** SET ac_bal = ac_bal - line_cost WHERE ac_number = 1234; (The third INVOICE has non-num data in ac_bal so the SQL Statement fails) COMMIT; (this commits the entire transaction) So the Journal will have • BGIN BFOR + AFTR (for ACCOUNT) BFOR + AFTER + BFOR + AFTR (for 1st and 2nd Invoices) RTSV (Rollback to the BFOR of 1st Invoice not the BGIN) COMT (Will commit all from BGIN, except those RSTV) • ENDJ Testing has taken a lot of time: IDMS handles this Jrnl record perfectlty. Rollbaclk and RFWD Some other utilities need amending to cope. ca IUA

Java Access to CA-IDMS Data at BT **Two Phase Commit**

- New requirement to coordinate with another DBMS (not mainframe)
- 2PC Generates still more Jrnl record types

DIND: Distributed transaction - In Doubt

DCOM: Co-ordinator Says COMMIT

DBAK: Co-ordinator says Back out

DPND: Co-ordinator Ask to wait before final Forget

DFGT: Co-ordinators says All done

Untill the Final DFGT, Transactions can be left In Doubt and not complete and locks still held

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Java Access to CA-IDMS Data at BT **Two Phase Commit**

New requirement to coordinate with another DBMS (not mainframe)

Testing journal processing Very difficult

Need to generate each scenario with various 2PC rollback and **Commit situations**

Need to coordinate

with other DBMS DBAs

and separately the 2PC Trasnaction manager.

Invoked from Java

Then SQL to IDMS and SQL to OtherDB

We are looking at Atomikos as a XM.





Java Access to CA-IDMS Data at BT Schema definitions

- Separate SQL Schemas to be defined
 - CREATE SCHEMA SQL SCCS00 FOR NETWORK SCHEMA SCCS00 v 292;

Exposes all IDMS records as Tables

Cannot be used with Table procedures

CREATE SCHEMA SQL_SCCS00_FK FOR NETWORK SCHEMA SCCS00 v 292 WITH VIRTUAL KEYS;

Exposes all IDMS records as Tables PLUS exposes all Set Owners Rowids/DBKEYS

Cannot be used with Table procedures

Used for new UPDATE Statements

CREATE SCHEMA SQLCSS00 FOR SQL;

Used to hang all Table Procedures, Keys, Views etc.

EG CREATE table Procedure SQLCSS00.VAL-IT

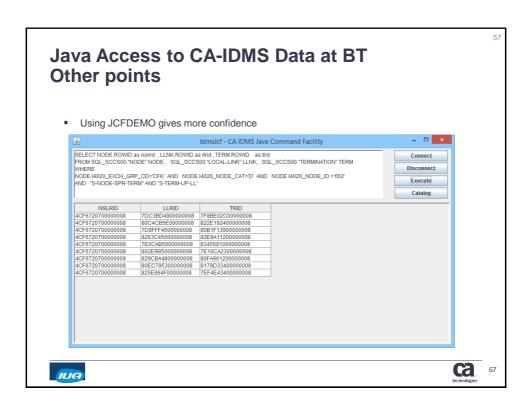
Used to define all other SQL entities used.

CREATE SCHEMA SQL_SCCS00 FOR NETWORK SCHEMA SCCS00IX v 292;

Used to satisfy reading non-MA system Indexes



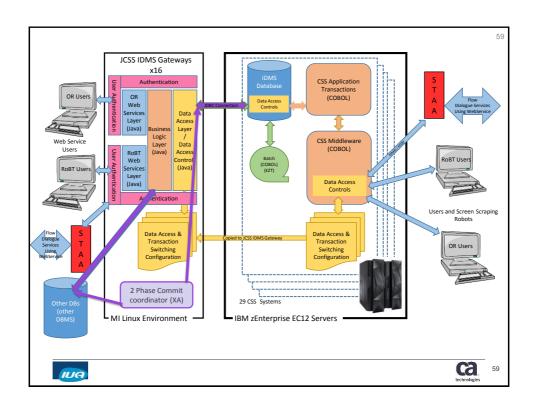




Java Access to CA-IDMS Data at BT Summary BT are continuing to exploit the modernisation techniques provided in IDMS to great effect CA continue to enhance and exploit the SQL and Java drivers Update SQL is coming 2 phase Commit processing is coming One last thing... Remember that diagram?

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